

Declarative Programming and (Co)Induction

Prolog exercises, part 2

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Easy exercises

1. try out the following non ground goals, with the predicates defined in the previous lab experience. For instance

```
?- is_nat(N).  
?- odd(N).  
?- even(N).  
?- div(s(s(s(z))), N).  
?- div(N, s(s(s(s(s(z)))))).
```

Consider both inductive and coinductive cases.

2. Define the predicate *add/3* s.t. *add*(t_1, t_2, t_3) holds iff t_1, t_2 , and t_3 are natural numbers and $t_3 = t_1 + t_2$.
Try out the goal *?- add*($N, M, s(s(z))$) with both the inductive and coinductive interpretations.

Less easy exercise

1. Implement the typechecking rules of the simply typed lambda-calculus as defined on pag.52 in the slides on functional programming.

Hints: Define the predicate *typeof/2* for ground terms (that is, where the type environment is implicitly empty), based on the auxiliary predicate *typeof/3* that takes also a type environment.

To implement the type environment you may use the library *assoc* (with `:- use_module(library(assoc)).`) and then the three predicates *empty_assoc/1* (to return an empty environment), *get_assoc/3* (to check the type of a variable), and *put_assoc/4* to update an environment (see the on-line documentation at <http://www.swi-prolog.org/pldoc/refman/>).

For representing the terms of the language, see the suggested syntax in the queries below.

```
?- E = fun(x : bool -> x), typeof(E, RT).  
?- E = fun(x : T -> x), typeof(E, RT).  
?- E = fun(f1 : T1 -> fun(f2 : T2 -> fun(x : T -> app(f1, app(f2, x))))), typeof(E, RT).  
?- E = fun(x : T -> app(x, x)), typeof(E, RT).  
?- E = fun(x : T -> app(x, x)), typeof(app(E, E), RT).  
?- E = fun(x : T -> app(x, x)), typeof(app(app(E, E), true), false), RT).  
?- E = fun(x : X -> fun(y : Y -> if(x, y, x))), typeof(E, RT).  
?- E1 = fun(x : X -> app(f, app(x, x))), E = fun(f : F -> app(E1, E1)), typeof(E, RT).  
?- F = fun(x : T -> x), E = fun(f : FT -> if(true, app(f, true), app(f, false))), typeof(E, RT).  
?- F = fun(x : T -> x), E = fun(f : FT -> if(true, app(f, true), app(f, F))), typeof(E, RT).
```

Try out the queries, and motivate the computed answers.

Very challenging exercise

- (a) Rational numbers in the interval $[0, 1]$ can be represented by regular lists of digits. For instance the list L s.t. $L = [1|L]$ corresponds to the repeating decimal $0.\hat{1}$, that is, $\frac{1}{9}$ (by the way, recall that $0.\hat{9} = 1$, see <http://en.wikipedia.org/wiki/0.999...>). Note also that finite decimal numbers can be seen as a specific case of repeating decimals; for instance, $0.1 = 0.1\hat{0}$.

Using regular lists to represent decimals in $[0, 1]$, implements the predicate *addrat/3* s.t. *addrat*(t_1, t_2, t_3) holds iff t_1, t_2 , and t_3 are repeating decimals and $t_1 + t_2 = t_3$.

Hints: use built-in numbers to represent digits, and built-in predicates to compute addition, and integer division with remainder (example X is $3 + 5$, Y is $5 // 10$, Z is $5 \bmod 10$).

Use an auxiliary predicates that computes also the regular list of all carries for all positions.